

# Graph Coloring Game

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Krzysztof Pióro

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Theoretical Computer Science

## Graph coloring game

In the *graph coloring game* two players **Alice** and **Bob** are given graph **G** and a set of **k** colors. **Alice** and **Bob** take turns, **coloring properly** an uncolored vertex.

- **Alice** wins when graph is completely colored
- **Bob** wins otherwise

# Definition

## Graph coloring game

In the *graph coloring game* two players **Alice** and **Bob** are given graph **G** and a set of **k** colors. **Alice** and **Bob** take turns, **coloring properly** an uncolored vertex.

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## Game chromatic number

The game chromatic number of a graph **G**, denoted by  $\chi_g(\mathbf{G})$ , is the minimum number of colors needed for **Alice** to win the graph coloring game on **G**.

## Example

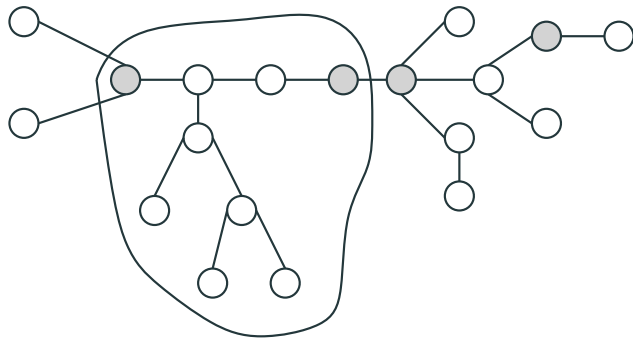


**Upper bound [Faigle, Kern, Kierstead, Trotter 1993]**

If  $\mathbf{G}$  is a tree, then  $\chi_g(\mathbf{G}) \leq 4$ .

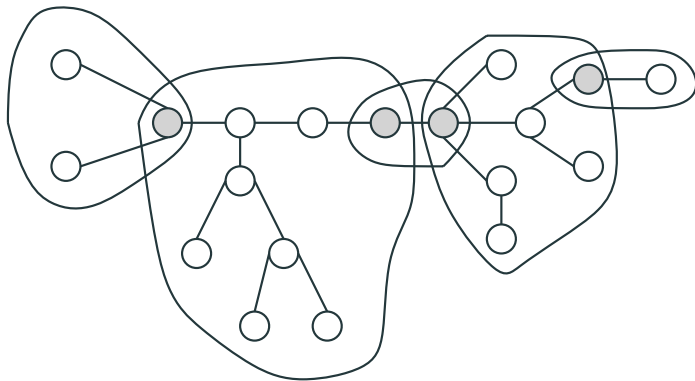
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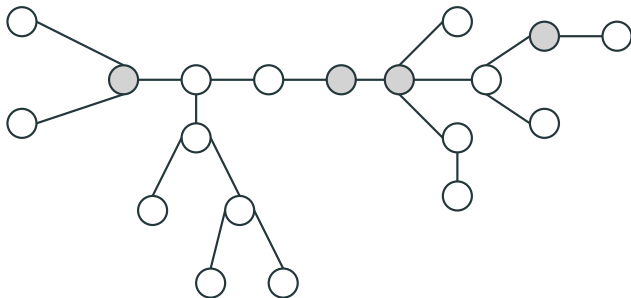
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# Trees

Upper bound [Faigle, Kern, Kierstead, Trotter 1993]

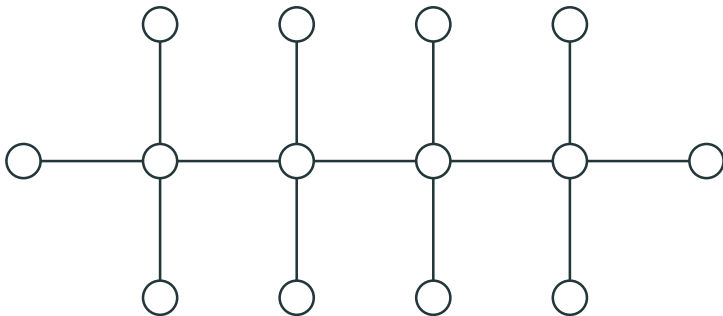
If  $\mathbf{G}$  is a tree, then  $\chi_g(\mathbf{G}) \leq 4$ .





## Lower bound [Bodlaender 1991]

There exists a tree  $\mathbf{T}$ , such that  $\chi_g(\mathbf{T}) = 4$



# Game coloring number

## Marking game

In the *marking game* two players **Alice** and **Bob** are given graph **G**. During the game they create linear order **L** of vertices of graph **G**. They alternate turns with **Alice** playing first. In each move player selects a vertex from the remaining vertices and puts it at the end of **L**.

- **Alice**'s goal is to **minimize** back degree of **L**
- **Bob**'s goal is to **maximize** back degree of **L**

# Game coloring number

## Marking game

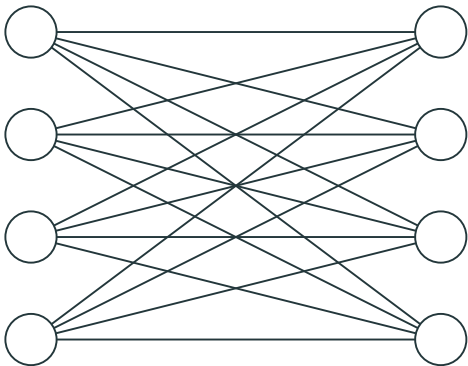
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- **Alice**'s goal is to **minimize** back degree of **L**
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## Game coloring number

The game coloring number of a graph **G**, denoted by  $\text{col}_g(\mathbf{G})$ , is equal to  $k + 1$ , where  $k$  is back degree of a linear order **L**, which is produced by playing the **marking game** with both players using their **optimal** strategies.

## Game coloring number



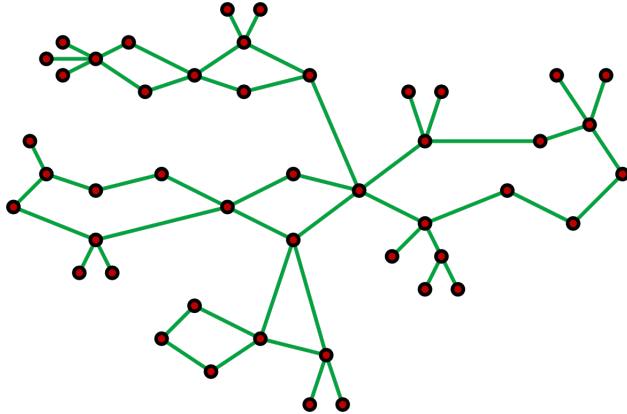
$$\text{col}_g(\mathbf{G}) = n + 1$$

$$\chi_g(\mathbf{G}) = 3$$

### Lemma

Suppose  $\mathbf{G} = (\mathbf{V}, \mathbf{E})$  and  $\mathbf{E} = \mathbf{E}_1 \cup \mathbf{E}_2$ . Let  $\mathbf{G}_1 = (\mathbf{V}, \mathbf{E}_1)$  and  $\mathbf{G}_2 = (\mathbf{V}, \mathbf{E}_2)$ .  
Then  $\text{col}_g(\mathbf{G}) \leq \text{col}_g(\mathbf{G}_1) + \Delta(\mathbf{G}_2)$ .

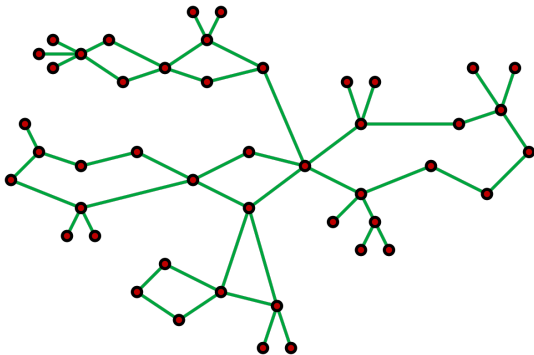
# Cactuses



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## Lemma

If  $G$  is a cactus then there is a matching  $M$  such that  $K - M$  is an acyclic graph.



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Thus edges of cactus  $\mathbf{G}$  can be splitted into two graphs  $\mathbf{G}_1, \mathbf{G}_2$ ,  
such that  $\mathbf{G}_1$  is a forest and  $\Delta(\mathbf{G}_2) \leq 1$ .



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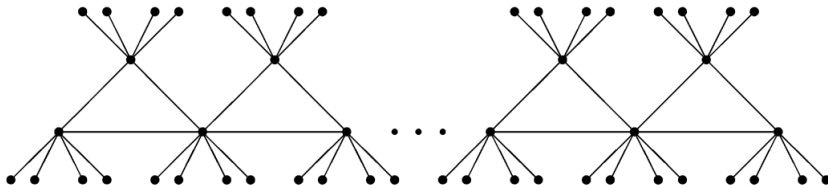
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## Upper bound [Sidorowicz 2006]

If  $\mathbf{G}$  is a cactus, then  $\chi_g(\mathbf{G}) \leq 5$ .

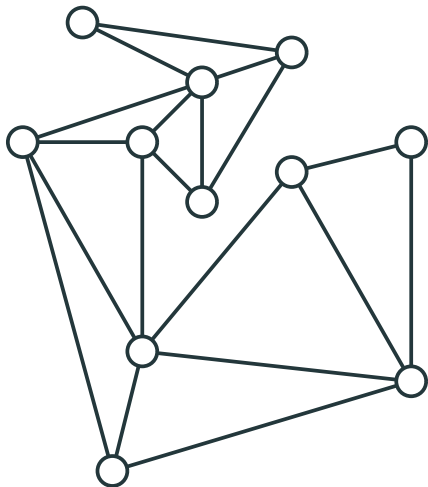
## Lower bound [Sidorowicz 2006]

There exists a cactus  $\mathbf{K}$ , such that  $\chi_g(\mathbf{K}) = 5$



$G$  contains 7 triangles

## Outerplanar graphs



## Theorem

For every maximal outerplanar graph there exists ordering of vertices  $\{v_1, v_2, \dots, v_n\}$  such that:

- $v_1v_2$  is an edge on the outer face,
- $\forall_{i>2}$   $v_i$  has exactly 2 neighbors on left. Let  $i_1 < i_2$  be the indices of these neighbors.

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## Properties

- $v_{i_1}$  is adjacent to  $v_{i_2}$
- $i \neq j \Rightarrow \{i_1, i_2\} \neq \{j_1, j_2\}$

## Lemma

For any vertex  $v_k$ , there are at most two vertices  $v_i, v_j$  such that  $i_2 = j_2 = k$ .

## Collary

Edges of every outerplanar graph can be splitted into two trees  $\mathbf{T}_1, \mathbf{T}_2$  such that  $\Delta(\mathbf{T}_2) \leq 3$ ,

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If  $G$  is an outerplanar graph, then  $\chi_g(G) \leq 7$ .

## Lower bound [Kierstead, Trotter 1994]

There exists an outerplanar graph  $T$ , such that  $\chi_g(T) = 6$

# Results

Class	Lower bound	Upper bound
Forests	4	4
Cactuses	5	5
Outerplanar graphs	6	7
Planar graphs	7	17
Interval graphs	$2\omega$	$3\omega - 2$

# Open problems

## Open problem 1

Suppose **Alice** has a winning strategy for the vertex coloring game on a graph **G** with  $k$  colors. Does she have one for  $k + 1$  colors?

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## Open problem 2

Is there a function  $f$  such that, if **Alice** has a winning strategy for the vertex coloring game on a graph  $\mathbf{G}$  with  $k$  colors, then Alice has a winning strategy on  $\mathbf{G}$  with  $f(k)$ .

# References

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