# A note on polynomials and $f$-factors of graphs 

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Graph k-factor: a spanning subgraph that is also k -regular.

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Graph k -factor: a spanning subgraph that is also k -regular.
Examples:

- 1-factor of a graph is it's perfect matching.
- 2-factor of a graph is it's cycle cover.


## $f$-factors

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To define what $f$-factor is we assume that we are given graph $G$ and a function $f: V(G) \rightarrow 2^{\mathbb{N}}$ such that:

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\forall_{v \in V(G)} f(v) \subseteq\{0, \ldots, \operatorname{deg}(v)\}
$$

Now we define $f$-factor as $G$ 's spanning subgraph $H$ such that:

$$
\forall_{v \in V(H)} \operatorname{deg}_{H}(v) \in f(v)
$$

## Conditions for $f$-factor existence

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In a case $\forall_{v \in V}|f(v)|=1$ we can show that existence of $f$-factor is equivalent to existence of perfect matching in some graph $G^{\prime}$.

Unfortulately, if $|f(v)|>1$ is allowed, then we do not expect to find an elegant condition that is both necessary and sufficient.
Even if $\forall_{v \in V}|f(v)| \in\{1,2\}$, then (for some $f$ s) deciding if there is $f$-factor in graph $G$ is NP-complete by edge 3-coloring reduction (The factorization of graphs. II. L. LOVASZ).

## Results from this paper

We cannot easily solve the decision problem, but we might look for conditions that are just sufficient. Authors give a following rule:

## Theorem

Let $G=(V, E)$ be a graph and suppose that $f$ satisfies:

$$
|f(v)|>\lceil\operatorname{deg}(v) / 2\rceil
$$

for every $v \in V$. Then $G$ has an $f$-factor.

## Results from this paper p. 2

To give a second theorem we need to first define a partial $f$-factor, which is $\bar{f}$-factor of the same graph with $\bar{f}(v)=f(v) \cup\{0\}$. We say that partial $f$-factor is non-trivial if it is not empty (there is at least one egde in the corresponding spanning subgraph).

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We say that partial $f$-factor is non-trivial if it is not empty (there is at least one egde in the corresponding spanning subgraph).

Now we can write down a statement of the second theorem from the paper:

## Theorem

Let $G=(V, E)$ be a graph, and let $f$ satisfy

$$
|E|>\sum_{v \in V}\left|f(v)^{c}-\{0\}\right|
$$

where $f(v)^{c}=\{0,1, \ldots, \operatorname{deg}(v)\}-f(v)$. Then $G$ contains a non-trivial partial $f$-factor.

## Combinatrial Nullstellensatz

Proofs of both theorems will take advantage of Combinatorial Nullstellensatz:

## Theorem

(Combinatorial Nullstellensatz) Let $g \in \mathbb{F}\left[X_{1}, X_{2}, \ldots, X_{n}\right]$ be a polynomial, and suppose the coefficient of the monomial $\prod_{i=1}^{n} X_{i}^{t_{i}}$ in $g$ is non-zero, where $t_{1}+\ldots+t_{n}$ is the total degree of $g$. Then, for any sets $S_{1}, \ldots, S_{n} \subset \mathbb{F}$ with $\left|S_{1}\right|>t,\left|S_{2}\right|, \ldots,\left|S_{n}\right|>t_{n}$, there exists $x \in S_{1} \times \ldots \times S_{n}$ such that $g(x) \neq 0$.

## First theorem: proof

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for every $v \in V$. Then $G$ has an $f$-factor.

First we define a polynomial:

$$
g=\prod_{v \in V} \prod_{c \in f(v)^{c}}\left(\sum_{e \ni v} X_{e}-c\right)
$$

Let's take $S_{1}=S_{2}=\ldots=S_{n}=\{0,1\}$, and look at the value of $X_{e}$ as a choice of including this edge into our subgraph.

## First theorem: proof

$$
g=\prod_{v \in V} \prod_{c \in f(v)^{c}}\left(\sum_{e \ni v} X_{e}-c\right)
$$

Now we can interpret existence of $x \in S_{1} \times \ldots \times S_{n}$ such that $g(x) \neq 0$ as existence of the $f$-factor: check if we can choose edges for our subgraph $H$ in a way such that there is no vertex $v$ for which $\operatorname{deg}_{H}(v) \in f(v)^{c}$.

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So now we need to prove that in $g$ we have a monomial of the form

$$
a \prod_{e \in E} X_{e}^{t_{e}}, a \neq 0
$$

where

$$
\forall_{e \in E} t_{e} \in\{0,1\}, \sum_{e \in E} t_{e}=\sum_{v \in V}\left|f(v)^{c}\right|
$$

## First theorem: proof

To achieve this we want to show that there is a function $R: V \rightarrow 2^{\mathbb{N}}$, such that:

$$
\forall_{v \in v}|R(v)|=\left|f(v)^{c}\right|
$$

and

$$
u, v \in V \wedge u \neq v \Longrightarrow R(u) \cap R(v)=\emptyset
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Now we can choose a monomial

$$
\prod_{v \in V} \prod_{e \in R(v)} X_{e}
$$

and it will have desired properties.

## First theorem: proof

How to construct such $R$ ? We can prove that there is an orientation of edges such that every vertex $v$ has outdegree at least $\left\lfloor\frac{1}{2} \operatorname{deg}(v)\right\rfloor$.

A simple argument is that if we add a new vertex, which is connected to vertices with odd degree, then we can orient edges as in a traverse of an Eulerian cycle.

Now, since we have $|f(v)|>\lceil\operatorname{deg}(v) / 2\rceil$, we know that $\left|f(v)^{c}\right| \leq\left\lfloor\frac{1}{2} \operatorname{deg}(v)\right\rfloor$, so $R$ can be constructed.

## Second theorem: proof

## Theorem

Let $G=(V, E)$ be a graph, and let $f$ satisfy

$$
|E|>\sum_{v \in V}\left|f(v)^{c}-\{0\}\right|
$$

where $f(v)^{c}=\{0,1, \ldots, \operatorname{deg}(v)\}-f(v)$. Then $G$ contains a non-trivial partial $f$-factor.

Now our polynomial will be:

$$
g=\prod_{v \in V} \prod_{c \in f(v)^{c}-\{0\}}\left(1-\frac{\sum_{e \ni v} X_{e}}{c}\right)-\prod_{e \in E}\left(1-X_{e}\right)
$$

## Second theorem: proof

$$
g=\prod_{v \in V} \prod_{c \in f(v)^{c}-\{0\}}\left(1-\frac{\sum_{e \ni v} x_{e}}{c}\right)-\prod_{e \in E}\left(1-X_{e}\right)
$$

Again we take $S_{0}=\ldots=S_{\mid} E \mid=\{0,1\}$. Now if there exists $x \neq 0$ such that $g(x) \neq 0$ then for this $x$ we have that

$$
\prod_{v \in V} \prod_{c \in f(v)^{c}-\{0\}}\left(1-\frac{\sum_{e \ni v} x_{e}}{c}\right) \neq 0
$$

and so

$$
\sum_{v \ni e} x_{e} \in f(x) \cup\{0\}
$$



## Second theorem: proof

$$
g=\prod_{v \in V} \prod_{c \in f(v)^{c}-\{0\}}\left(1-\frac{\sum_{e \ni v} X_{e}}{c}\right)-\prod_{e \in E}\left(1-X_{e}\right)
$$

Now we need to show that largest degree monomial fulfils assumptions of Combinatorial Nullstellensatz.
We know that for this graph

$$
\sum_{v \in V}\left|f(v)^{c}-\{0\}\right|<|E|
$$

So it's obvious that the largest monomial is just

$$
(-1)^{|E|+1} \prod_{e \in E} X_{e}
$$

