Clustered Coloring of Graphs Excluding a Subgraph and a Minor
Chun-Hung Liu, David R. Wood
[2019+]

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& \chi(G) \leqslant C \cdot t \cdot(\log \log t)^{6}
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[Norin,Song,Postle 2020]
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\forall_{s, t, H} \exists_{\eta=\eta(s, t, H)} \forall_{G}
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$$
G \text { has no } H \text {-minor }
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G \text { has no } K_{s, t} \text {-subgraph }
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2-clustered 3 coloring

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# Technical statement: (optimal) 

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$\chi_{c}\left(\mathcal{G}_{s}\right) \leqslant s$
$\chi_{c}\left(\mathcal{G}_{s}\right) \leqslant s+2$
$\operatorname{tw}(G) \leqslant \omega$

$G$ has-no $H$-minor
$G$ has no $K_{s, t}$-subgraph
$(s+1)$
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$(s+1)$ choosable
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$(A, B)$ is a Separation if $A \cup B=G$ and $E(A) \cap E(B)=\emptyset$. $\operatorname{ord}(A, B):=|V(A \cap B)|$

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(T1) $\operatorname{ord}(A, B)<\theta$, either $(A, B) \in \mathcal{T}$ or $(B, A) \in \mathcal{T}$
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(T3) $(A, B) \in \mathcal{T}$, then $V(A) \neq V(G)$
Example: $C$ fixed cycle in $G$
$\mathcal{T}=\{(A, B):$ ord $=1, C \subset B\}$

$$
\forall_{s, t, \omega} \exists_{\eta=\eta(s, t, \omega)} \forall_{G}
$$

- $\operatorname{tw}(G) \leqslant \omega$
- $G$ has no $K_{s, t}$-subgraph
$\Rightarrow G$ is $(s+1)$-choosable with clustering $\eta$
$(A, B)$ is a Separation if $A \cup B=G$ and $E(A) \cap E(B)=\emptyset$.
$\operatorname{ord}(A, B):=|V(A \cap B)|$


B
(T1) $\operatorname{ord}(A, B)<\theta$, either $(A, B) \in \mathcal{T}$ or $(B, A) \in \mathcal{T}$
(T2) $\left(A_{1}, B_{1}\right),\left(A_{2}, B_{2}\right),\left(A_{3}, B_{3}\right) \in \mathcal{T}$, then $A_{1} \cup A_{2} \cup A_{3} \neq G$
(T3) $(A, B) \in \mathcal{T}$, then $V(A) \neq V(G)$
Example: $C$ fixed cycle in $G$
$\mathcal{T}=\{(A, B):$ ord $=1, C \subset B\}$
$\mathcal{T}$ - set of some separation of order $<\theta$ $\mathcal{T}$ is a Tangle of order $\theta$ if:

$$
\forall_{s, t, \omega} \exists_{\eta=\eta(s, t, \omega)} \forall_{G}
$$

- $\operatorname{tw}(G) \leqslant \omega$
- $G$ has no $K_{s, t}$-subgraph
$\Rightarrow G$ is $(s+1)$-choosable with clustering $\eta$
$(A, B)$ is a Separation if $A \cup B=G$ and $E(A) \cap E(B)=\emptyset$.
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B
(T1) $\operatorname{ord}(A, B)<\theta$, either $(A, B) \in \mathcal{T}$ or $(B, A) \in \mathcal{T}$
(T2) $\left(A_{1}, B_{1}\right),\left(A_{2}, B_{2}\right),\left(A_{3}, B_{3}\right) \in \mathcal{T}$, then $A_{1} \cup A_{2} \cup A_{3} \neq G$
(T3) $(A, B) \in \mathcal{T}$, then $V(A) \neq V(G)$
Example: $C$ fixed cycle in $G$
$\mathcal{T}=\{(A, B):$ ord $=1, C \subset B\}$
$\mathcal{T}$ - tangle of order 2

$$
\forall_{s, t, \omega} \exists_{\eta=\eta(s, t, \omega)} \forall_{G}
$$

- $\operatorname{tw}(G) \leqslant \omega$
- $G$ has no $K_{s, t}$-subgraph
$\Rightarrow G$ is $(s+1)$-choosable with clustering $\eta$

Advanced example: $G=\boxplus_{k}$ $\mathcal{T}=\{(A, B):$ ord $<k$, full row $\subset B\}$
$(A, B)$ is a Separation if $A \cup B=G$ and $E(A) \cap E(B)=\emptyset$.
$\operatorname{ord}(A, B):=|V(A \cap B)|$
$\mathcal{T}$ - set of some separation of order $<\theta$


B $\mathcal{T}$ is a Tangle of order $\theta$ if:
(T1) $\operatorname{ord}(A, B)<\theta$, either $(A, B) \in \mathcal{T}$ or $(B, A) \in \mathcal{T}$
(T2) $\left(A_{1}, B_{1}\right),\left(A_{2}, B_{2}\right),\left(A_{3}, B_{3}\right) \in \mathcal{T}$, then $A_{1} \cup A_{2} \cup A_{3} \neq G$
(T3) $(A, B) \in \mathcal{T}$, then $V(A) \neq V(G)$
Example: $C$ fixed cycle in $G$
$\mathcal{T}=\{(A, B):$ ord $=1, C \subset B\}$
$\mathcal{T}$ - tangle of order 2

$$
\begin{aligned}
& \forall_{s, t, \omega} \exists_{\eta=\eta(s, t, \omega)} \forall_{G} \\
& \text { - } \operatorname{tw}(G) \leqslant \omega \\
& \text { - } G \text { has no } K_{s, t} \text {-subgraph } \\
& :=|V(A \cap B)| \\
& \Rightarrow G \text { is }(s+1) \text {-choosable with clustering } \eta \\
& \mathcal{T} \text { - ...order }<\theta \\
& \text { (T1) }(A, B) \in \mathcal{T} \text { or }(B, A) \in \mathcal{T} \\
& \text { (T2) } A_{1} \cup A_{2} \cup A_{3} \neq G \\
& \text { (T3) } V(A) \neq V(G)
\end{aligned}
$$

$$
\forall_{s, t, \omega} \exists_{\eta=\eta(s, t, \omega)} \forall_{G}
$$

- $\operatorname{tw}(G) \leqslant \omega$
- $G$ has no $K_{s, t}$-subgraph

$$
:=|V(A \cap B)|
$$

$\Rightarrow G$ is $(s+1)$-choosable with clustering $\eta$

- $\operatorname{tw}(G) \leqslant \omega$
$\mathcal{T}$ - ...order $<\theta$
(T1) $(A, B) \in \mathcal{T}$ or $(B, A) \in \mathcal{T}$
(T2) $A_{1} \cup A_{2} \cup A_{3} \neq G$
(T3) $V(A) \neq V(G)$

$$
:=\text { no tangle of order } \omega+2
$$

$$
\begin{aligned}
& \forall_{s, t, \omega} \exists_{\eta=\eta(s, t, \omega)} \forall_{G} \\
& \text { - } \operatorname{tw}(G) \leqslant \omega \quad(:=\text { no tangle of or der } \omega+2) \\
& \text { - } G \text { has no } K_{s, t} \text {-subgraph } \\
& \Rightarrow G \text { is }(s+1) \text {-choosable with clustering } \eta \\
& \operatorname{ord}(A, B) \\
& :=|V(A \cap B)| \\
& \mathcal{T} \text { - ...order }<\theta \\
& \text { (T1) }(A, B) \in \mathcal{T} \text { or }(B, A) \in \mathcal{T} \\
& \text { (T2) } A_{1} \cup A_{2} \cup A_{3} \neq G \\
& \text { (T3) } V(A) \neq V(G)
\end{aligned}
$$

$$
\begin{aligned}
& \forall_{s, t, \omega} \exists_{\eta=\eta(s, t, \omega)} \forall_{G} \quad \bullet \operatorname{tw}(G) \leqslant \omega \quad(:=\text { no tangle of or oder } \omega+2) \\
& \Rightarrow G \text { is }(s+1) \text {-choosable with clustering } \eta \\
& \mathcal{T} \text { - ...order }<\theta \\
& \text { (T1) }(A, B) \in \mathcal{T} \text { or }(B, A) \in \mathcal{T} \\
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\end{aligned}
$$

$$
\begin{aligned}
& \forall_{s, t, \omega} \exists_{\eta=\eta(s, t, \omega)} \forall_{G} \quad \bullet \operatorname{tw}(G) \leqslant \omega \quad(:=\text { no tangle of or der } \omega+2) \\
& \text { - } G \text { has no } K_{s, t} \text {-subgraph } \\
& \Rightarrow G \text { is }(s+1) \text {-choosable with clustering } \eta \\
& \text { (T1) }(A, B) \in \mathcal{T} \text { or }(B, A) \in \mathcal{T} \\
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& \text { (T3) } V(A) \neq V(G)
\end{aligned}
$$



$$
\begin{aligned}
& \forall_{s, t, \omega} \exists_{\eta=\eta(s, t, \omega)} \forall_{G} \quad \bullet \operatorname{tw}(G) \leqslant \omega \quad(:=\text { no tangle of or der } \omega+2) \\
& \text { - } G \text { has no } K_{s, t} \text {-subgraph } \\
& \Rightarrow G \text { is }(s+1) \text {-choosable with clustering } \eta \\
& \operatorname{ord}(A, B) \\
& :=|V(A \cap B)| \\
& \mathcal{T} \text { - ...order }<\theta \\
& \text { (T1) }(A, B) \in \mathcal{T} \text { or }(B, A) \in \mathcal{T} \\
& \text { (T2) } A_{1} \cup A_{2} \cup A_{3} \neq G \\
& \text { (T3) } V(A) \neq V(G)
\end{aligned}
$$



$$
\begin{aligned}
& \forall_{s, t, \omega} \exists_{\eta=\eta(s, t, \omega)} \forall_{G} \\
& \text { - } \operatorname{tw}(G) \leqslant \omega \quad(:=\text { no tangle of order } \omega+2) \\
& \Rightarrow G \text { is }(s+1) \text {-choosable with clustering } \eta \\
& X \subset V(G) \\
& \text { ord }(A, B) \\
& :=|V(A \cap B)| \\
& \text { (T1) }(A, B) \in \mathcal{T} \text { or }(B, A) \in \mathcal{T} \\
& \text { (T2) } A_{1} \cup A_{2} \cup A_{3} \neq G \\
& \text { (T3) } V(A) \neq V(G)
\end{aligned}
$$

$$
\left|N^{\geqslant s}(X)\right| \leqslant(t-1) \cdot\binom{|X|}{s}+1
$$

$$
\forall_{s, t, \omega} \exists_{\eta=\eta(s, t, \omega)} \forall_{G}
$$

- $\operatorname{tw}(G) \leqslant \omega \quad(:=$ no tangle of of der $\omega+2)$
$\Rightarrow G$ is $(s+1)$-choosable with clustering $\eta$
We start with $|L(v)|=s+1$
$\operatorname{ord}(A, B)$

$$
:=|V(A \cap B)|
$$

$\mathcal{T}$ - ...order $<\theta$
(T1) $(A, B) \in \mathcal{T}$ or $(B, A) \in \mathcal{T}$
(T2) $A_{1} \cup A_{2} \cup A_{3} \neq G$
(T3) $V(A) \neq V(G)$

$$
\left|N^{\geqslant s}(X)\right| \leqslant f(|X|, s, t)
$$

$\forall_{s, t, \omega} \exists_{\eta=\eta(s, t, \omega)} \forall_{G}$

- $\operatorname{tw}(G) \leqslant \omega \quad(:=$ no tangle of order $\omega+2)$
$\Rightarrow G$ is $(s+1)$-choosable with clustering $\eta$
We start with $|L(v)|=s+1$
Iteratively we enlarge colored set $Y$ (until not too big)
$\operatorname{ord}(A, B)$

$$
:=|V(A \cap B)|
$$

$\mathcal{T}$ - ...order $<\theta$
(T1) $(A, B) \in \mathcal{T}$ or $(B, A) \in \mathcal{T}$
(T2) $A_{1} \cup A_{2} \cup A_{3} \neq G$
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$$
\left|N^{\geqslant s}(X)\right| \leqslant f(|X|, s, t)
$$

$\forall_{s, t, \omega} \exists_{\eta=\eta(s, t, \omega)} \forall_{G}$

- $\operatorname{tw}(G) \leqslant \omega \quad(:=$ no tangle of order $\omega+2)$
$\Rightarrow G$ is $(s+1)$-choosable with clustering $\eta$
We start with $|L(v)|=s+1$
Iteratively we enlarge colored set $Y$ (until not too big) Invariant:
- $G$ has no $K_{s, t}$-subgraph $\operatorname{ord}(A, B)$

$$
:=|V(A \cap B)|
$$

$\mathcal{T}$ - ...order $<\theta$
(T1) $(A, B) \in \mathcal{T}$ or $(B, A) \in \mathcal{T}$
(T2) $A_{1} \cup A_{2} \cup A_{3} \neq G$
(T3) $V(A) \neq V(G)$
$\left|N^{\geqslant s}(X)\right| \leqslant f(|X|, s, t)$
$\forall_{s, t, \omega} \exists_{\eta=\eta(s, t, \omega)} \forall_{G}$

- $\operatorname{tw}(G) \leqslant \omega \quad(:=$ no tangle of or der $\omega+2)$
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$$
\mathcal{T}-\ldots \text { order }<\theta
$$

(T1) $(A, B) \in \mathcal{T}$ or $(B, A) \in \mathcal{T}$
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$$
\left|N^{\geqslant s}(X)\right| \leqslant f(|X|, s, t)
$$

$\forall_{s, t, \omega} \exists_{\eta=\eta(s, t, \omega)} \forall_{G}$

- $\operatorname{tw}(G) \leqslant \omega \quad(:=$ no tangle of order $\omega+2)$
- $G$ has no $K_{s, t}$-subgraph $\operatorname{ord}(A, B)$

$$
:=|V(A \cap B)|
$$

$\Rightarrow G$ is $(s+1)$-choosable with clustering $\eta$
We start with $|L(v)|=s+1$
Iteratively we enlarge colored set $Y$ (until not too big) Invariant:

$$
|L(v)|=1 \quad N^{\geqslant s}(Y)
$$

removed nbs colors!

$$
|L(v)| \geqslant 2
$$

$N^{<s}(Y)$

$$
|L(v)|=s+1
$$

$\forall_{s, t, \omega} \exists_{\eta=\eta(s, t, \omega)} \forall_{G}$

- $\operatorname{tw}(G) \leqslant \omega \quad(:=$ no tangle of order $\omega+2)$
- $G$ has no $K_{s, t}$-subgraph $\operatorname{ord}(A, B)$

$$
:=|V(A \cap B)|
$$

$\Rightarrow G$ is $(s+1)$-choosable with clustering $\eta$
We start with $|L(v)|=s+1$
Iteratively we enlarge colored set $Y$ (until not too big)
Invariant:

$$
N^{\geqslant s}(Y) \quad \text { Start with } Y=\{v\} \text { and any color }
$$

removed nbs colors!

$$
|L(v)| \geqslant 2
$$

$N^{<s}(Y)$

$$
|L(v)|=s+1
$$

$\forall_{s, t, \omega} \exists_{\eta=\eta(s, t, \omega)} \forall_{G}$

- $\operatorname{tw}(G) \leqslant \omega \quad(:=$ no tangle of order $\omega+2)$
- $G$ has no $K_{s, t}$-subgraph $\operatorname{ord}(A, B)$

$$
:=|V(A \cap B)|
$$

$\Rightarrow G$ is $(s+1)$-choosable with clustering $\eta$
We start with $|L(v)|=s+1$
Iteratively we enlarge colored set $Y$ (until not too big)
Invariant:

$$
N^{\geqslant s}(Y)
$$

Start with $Y=\{v\}$ and any color
removed nbs colors!

$$
\text { If } N^{\geqslant s}(Y)=\emptyset \text { color some } z \in N^{<s}(Y)
$$

$$
|L(v)| \geqslant 2
$$

$$
|L(v)|=1
$$

$N^{<s}(Y)$

$$
|L(v)|=s+1
$$

$\forall_{s, t, \omega} \exists_{\eta=\eta(s, t, \omega)} \forall_{G}$

- $\operatorname{tw}(G) \leqslant \omega \quad(:=$ no tangle of order $\omega+2)$
$\Rightarrow G$ is $(s+1)$-choosable with clustering $\eta$
We start with $|L(v)|=s+1$
Iteratively we enlarge colored set $Y$ (until not too big)

$$
\begin{aligned}
& \operatorname{ord}(A, B) \\
& \quad:=|V(A \cap B)|
\end{aligned}
$$

$$
\mathcal{T}-\ldots \text { order }<\theta
$$

(T1) $(A, B) \in \mathcal{T}$ or $(B, A) \in \mathcal{T}$
(T2) $A_{1} \cup A_{2} \cup A_{3} \neq G$
(T3) $V(A) \neq V(G)$

$$
\left|N^{\geqslant s}(X)\right| \leqslant f(|X|, s, t)
$$

Start with $Y=\{v\}$ and any color If $N^{\geqslant s}(Y)=\emptyset$ color some $z \in N^{<s}(Y)$

## Otherwise:

$$
\forall_{s, t, \omega} \exists_{\eta=\eta(s, t, \omega)} \forall_{G}
$$

- $\operatorname{tw}(G) \leqslant \omega \quad(:=$ no tangle of or der $\omega+2)$
- $G$ has no $K_{s, t}$-subgraph $\operatorname{ord}(A, B)$

$$
:=|V(A \cap B)|
$$

$\Rightarrow G$ is $(s+1)$-choosable with clustering $\eta$
We start with $|L(v)|=s+1$
Iteratively we enlarge colored set $Y$ (until not too big)

$$
\mathcal{T}-\ldots \text { order }<\theta
$$

(T1) $(A, B) \in \mathcal{T}$ or $(B, A) \in \mathcal{T}$
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(T3) $V(A) \neq V(G)$

$$
\left|N^{\geqslant s}(X)\right| \leqslant f(|X|, s, t)
$$

Start with $Y=\{v\}$ and any color
If $N^{\geqslant s}(Y)=\emptyset$ color some $z \in N^{<s}(Y)$

## Otherwise:

$N^{<s}$
$\forall_{s, t, \omega} \exists_{\eta=\eta(s, t, \omega)} \forall_{G}$

- $\operatorname{tw}(G) \leqslant \omega \quad(:=$ no tangle of order $\omega+2)$
$\Rightarrow G$ is $(s+1)$-choosable with clustering $\eta$
We start with $|L(v)|=s+1$
Iteratively we enlarge colored set $Y$ (until not too big)


$$
\begin{aligned}
& \operatorname{ord}(A, B) \\
& \quad:=|V(A \cap B)|
\end{aligned}
$$

$$
\mathcal{T}-\ldots \text { order }<\theta
$$

(T1) $(A, B) \in \mathcal{T}$ or $(B, A) \in \mathcal{T}$
(T2) $A_{1} \cup A_{2} \cup A_{3} \neq G$
(T3) $V(A) \neq V(G)$

$$
\left|N^{\geqslant s}(X)\right| \leqslant f(|X|, s, t)
$$

Start with $Y=\{v\}$ and any color If $N^{\geqslant s}(Y)=\emptyset$ color some $z \in N^{<s}(Y)$

## Otherwise:

$$
\forall_{s, t, \omega} \exists_{\eta=\eta(s, t, \omega)} \forall_{G}
$$

- $\operatorname{tw}(G) \leqslant \omega \quad(:=$ no tangle of or der $\omega+2)$
$\Rightarrow G$ is $(s+1)$-choosable with clustering $\eta$

We start with $|L(v)|=s+1$
Iteratively we enlarge colored set $Y$ (until not too big)


$$
\begin{aligned}
& \operatorname{ord}(A, B) \\
& \quad:=|V(A \cap B)|
\end{aligned}
$$

- $G$ has no $K_{s, t}$-subgraph

$$
\mathcal{T}-\ldots \text { order }<\theta
$$

(T1) $(A, B) \in \mathcal{T}$ or $(B, A) \in \mathcal{T}$
(T2) $A_{1} \cup A_{2} \cup A_{3} \neq G$
(T3) $V(A) \neq V(G)$

$$
\left|N^{\geqslant s}(X)\right| \leqslant f(|X|, s, t)
$$

Start with $Y=\{v\}$ and any color
If $N^{\geqslant s}(Y)=\emptyset$ color some $z \in N^{<s}(Y)$

## Otherwise:

$$
\forall_{s, t, \omega} \exists_{\eta=\eta(s, t, \omega)} \forall_{G}
$$

- $\operatorname{tw}(G) \leqslant \omega \quad(:=$ no tangle of or der $\omega+2)$
$\Rightarrow G$ is $(s+1)$-choosable with clustering $\eta$
We start with $|L(v)|=s+1$
Iteratively we enlarge colored set $Y$ (until not too big)


$$
\begin{aligned}
& \operatorname{ord}(A, B) \\
& \quad:=|V(A \cap B)|
\end{aligned}
$$

$\mathcal{T}$ - ...order $<\theta$
(T1) $(A, B) \in \mathcal{T}$ or $(B, A) \in \mathcal{T}$
(T2) $A_{1} \cup A_{2} \cup A_{3} \neq G$
(T3) $V(A) \neq V(G)$

$$
\left|N^{\geqslant s}(X)\right| \leqslant f(|X|, s, t)
$$

Start with $Y=\{v\}$ and any color
If $N^{\geqslant s}(Y)=\emptyset$ color some $z \in N^{<s}(Y)$

## Otherwise:


$\forall_{s, t, \omega} \exists_{\eta=\eta(s, t, \omega)} \forall_{G}$

- $\operatorname{tw}(G) \leqslant \omega \quad(:=$ no tangle of order $\omega+2)$
- $G$ has no $K_{s, t}$-subgraph
$\operatorname{ord}(A, B)$

$$
:=|V(A \cap B)|
$$

$\Rightarrow G$ is $(s+1)$-choosable with clustering $\eta$
We start with $|L(v)|=s+1$
Iteratively we enlarge colored set $Y$ (until not too big)

$$
\mathcal{T}-\ldots \text { order }<\theta
$$

(T1) $(A, B) \in \mathcal{T}$ or $(B, A) \in \mathcal{T}$
(T2) $A_{1} \cup A_{2} \cup A_{3} \neq G$
(T3) $V(A) \neq V(G)$

$$
\left|N^{\geqslant s}(X)\right| \leqslant f(|X|, s, t)
$$

Start with $Y=\{v\}$ and any color
If $N^{\geqslant s}(Y)=\emptyset$ color some $z \in N^{<s}(Y)$
Otherwise:


$$
\begin{aligned}
& \forall_{s, t, \omega} \exists_{\eta=\eta(s, t, \omega)} \forall_{G} \\
& \text { - } \operatorname{tw}(G) \leqslant \omega \quad(:=\text { no tangle of order } \omega+2) \\
& \Rightarrow G \text { is }(s+1) \text {-choosable with clustering } \eta \\
& \text { We start with }|L(v)|=s+1 \\
& \text { Iteratively we enlarge colored set } Y \text { (until not too big) } \\
& \mathcal{T}_{\theta}:=\left\{(A, B)_{\theta}:|V(A) \cap Y| \leqslant 3 \theta\right\}|Y|>9 \theta \quad \theta:=\omega+2 \\
& \operatorname{ord}(A, B) \\
& :=|V(A \cap B)| \\
& \mathcal{T} \text { - ...order }<\theta \\
& \text { (T1) }(A, B) \in \mathcal{T} \text { or }(B, A) \in \mathcal{T} \\
& \text { (T2) } A_{1} \cup A_{2} \cup A_{3} \neq G \\
& \text { (T3) } V(A) \neq V(G) \\
& \left|N^{\geqslant s}(X)\right| \leqslant f(|X|, s, t)
\end{aligned}
$$

Start with $Y=\{v\}$ and any color
If $N^{\geqslant s}(Y)=\emptyset$ color some $z \in N^{<s}(Y)$
Otherwise:

$\forall_{s, t, \omega} \exists_{\eta=\eta(s, t, \omega)} \forall_{G}$

- $\operatorname{tw}(G) \leqslant \omega \quad(:=$ no tangle of order $\omega+2)$
- $G$ has no $K_{s, t}$-subgraph $\operatorname{ord}(A, B)$

$$
:=|V(A \cap B)|
$$

$\Rightarrow G$ is $(s+1)$-choosable with clustering $\eta$
We start with $|L(v)|=s+1$
Iteratively we enlarge colored set $Y$ (until not too big)
$\mathcal{T}_{\theta}:=\left\{(A, B)_{\theta}:|V(A) \cap Y| \leqslant 3 \theta\right\}|Y|>9 \theta \quad \theta:=\omega+2$
is not a tangle!
But:


Start with $Y=\{v\}$ and any color
If $N^{\geqslant s}(Y)=\emptyset$ color some $z \in N^{<s}(Y)$

Otherwise:


$$
\forall_{s, t, \omega} \exists_{\eta=\eta(s, t, \omega)} \forall_{G}
$$

- $\operatorname{tw}(G) \leqslant \omega \quad(:=$ no tangle of order $\omega+2)$
- $G$ has no $K_{s, t}$-subgraph $\operatorname{ord}(A, B)$

$$
:=|V(A \cap B)|
$$

$\Rightarrow G$ is $(s+1)$-choosable with clustering $\eta$
We start with $|L(v)|=s+1$
Iteratively we enlarge colored set $Y$ (until not too big)
$\mathcal{T}_{\theta}:=\left\{(A, B)_{\theta}:|V(A) \cap Y| \leqslant 3 \theta\right\}|Y|>9 \theta \quad \theta:=\omega+2$
$\mathcal{T}$ - ...order $<\theta$
(T1) $(A, B) \in \mathcal{T}$ or $(B, A) \in \mathcal{T}$
(T2) $A_{1} \cup A_{2} \cup A_{3} \neq G$
(T3) $V(A) \neq V(G)$

$$
\left|N^{\geqslant s}(X)\right| \leqslant f(|X|, s, t)
$$

is not a tangle!
But: (T2), (T3) - OK


Start with $Y=\{v\}$ and any color
If $N^{\geqslant s}(Y)=\emptyset$ color some $z \in N^{<s}(Y)$

Otherwise:


$$
\forall_{s, t, \omega} \exists_{\eta=\eta(s, t, \omega)} \forall_{G}
$$

- $\operatorname{tw}(G) \leqslant \omega \quad(:=$ no tangle of order $\omega+2)$
- $G$ has no $K_{s, t}$-subgraph $\operatorname{ord}(A, B)$

$$
:=|V(A \cap B)|
$$

$\Rightarrow G$ is $(s+1)$-choosable with clustering $\eta$
We start with $|L(v)|=s+1$
Iteratively we enlarge colored set $Y$ (until not too big)
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But: (T2), (T3) - OK $\Rightarrow \sim(\mathrm{T} 1)$ is not a tangle!


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## Hadwiger's clustered conjecture:

$\mathcal{G}_{s}=K_{s+1}$ minor free graphs $\chi_{c}\left(\mathcal{G}_{s}\right) \leqslant s$

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$K_{s, s,}$ has $K_{s+1}$ minor

## We proved



Technical statement: (optimal)

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\begin{aligned}
& (s+1) \text { choosable } \\
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& G \text { has no } K_{s, t} \text {-subgraph } \\
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source: Felix Reidl's website
- Results from $\sim 100$ pages

> culim 2.ses. long companion paper $\left(Y^{(i,-1, k+1)}-Y^{(i,-1, k)}\right) \cap X_{V\left(T_{t}\right)} \subseteq N_{G}\left[W_{0}^{(i,-1, k)}\right] \cap X_{V\left(T_{t}\right)} \subseteq N_{G}\left[\bigcup_{j^{\prime}=1}^{|V|-1} \bigcup_{S \in S_{j^{\prime}}} S\right] \subseteq \bigcup_{j^{\prime}=1}^{|\mathcal{L}|-1} T_{j^{\prime}} \subseteq \bigcup_{j^{\prime}=1}^{|\nu|-1} I_{f^{\prime}}$
So for every $k \in\left[0, w_{0}-1\right]$ and $q \in[0, s+1]$,
$\left(Y^{(i,-1, k, q+1)}-Y^{(i,-1, k, q)}\right) \cap I_{j} \cap X_{V\left(T_{t}\right)}-X_{t} \subseteq A_{L(k,-1, k, 1)}\left(Y_{1}^{(i,-1, k, k)} \cap \overline{I_{j}^{\bar{V}}}\right) \cap I_{j} \cap X_{V\left(T_{t}\right)}-X_{t}$ $\subseteq N_{G}^{\nabla_{0}^{*}}\left(Y^{(i,-1, k, k)} \cap \overline{T_{j}^{5}}\right) \cap I_{j} \cap X_{V\left(T_{t}\right)}-X_{t}$ $\subseteq N_{G}^{s s}\left(Y^{(i,-1, k, a)} \cap \overline{T_{j}^{o}} \cap X_{V\left(T_{t}\right)}\right)$.

