

On induced subgraphs with all degrees odd

by Alexander Scott

Presented by Jakub Siuta

January 05 2023

The even case

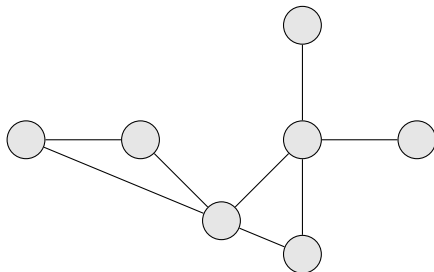
Gallai's Theorem

Every graph G has a vertex partition into two sets, each of which induces a subgraph with all degrees even

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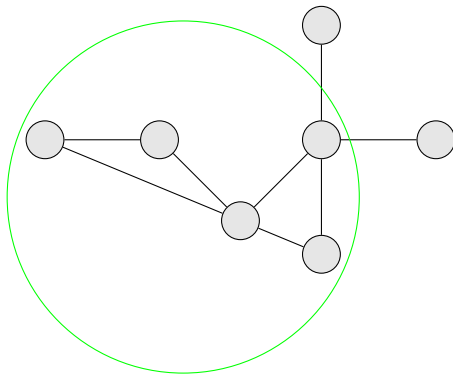
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- $f(n) \geq cn$ - A. Ferber, M. Krivelevich (2022)

Good Partitions

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Theorem

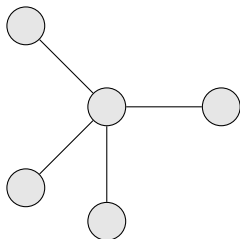
Let G be a graph. Then G has a good partition if and only if every component of G has even order.

Theorem proof

If G has a good partition, then it has to have even order, because of its vertices degree sum. To show the second implication, we will use induction.

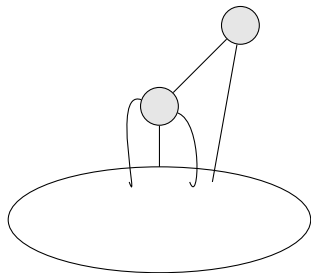
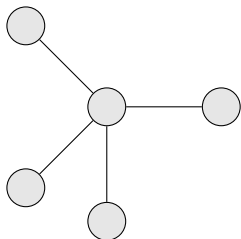
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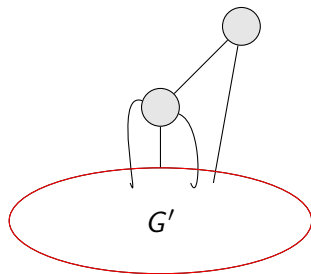
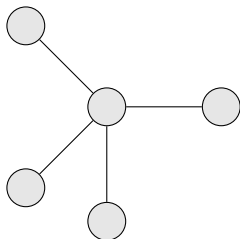
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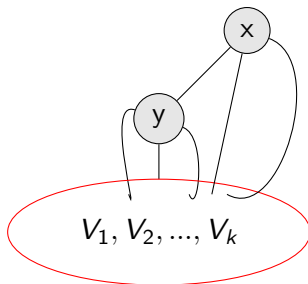


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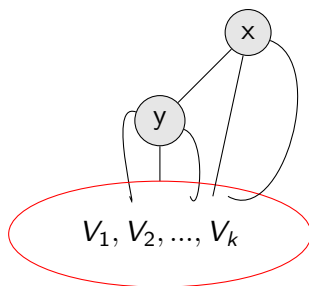
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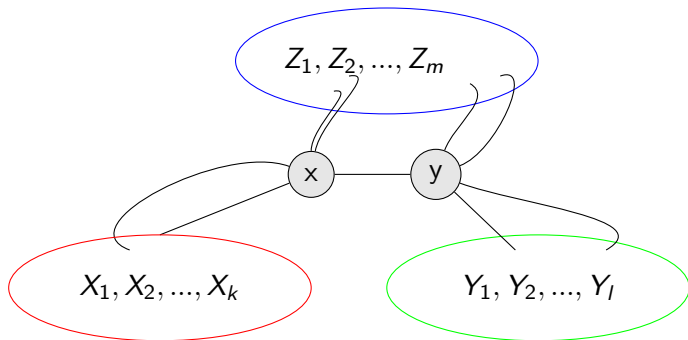


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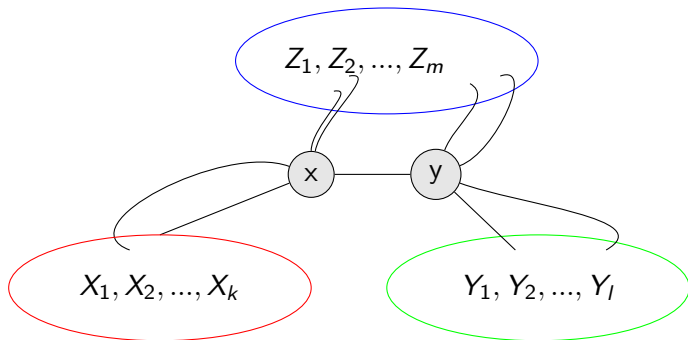


If at least one of V_i has even order, we can take G , and $G \setminus V_i$, if not, then all of V_i have odd order, so there is an even number of V_i

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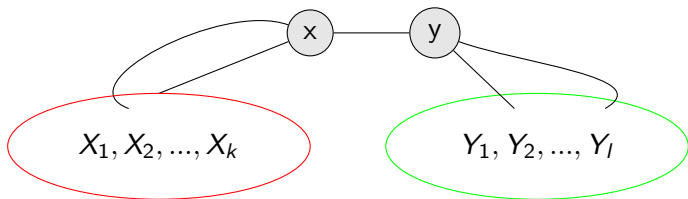


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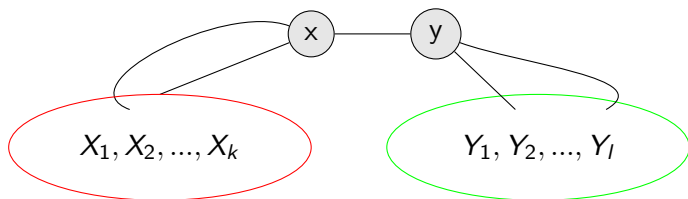


- $X = \{x\} \cup \bigcup_{i=1}^k V(X_i) \cup \bigcup_{i=2}^m V(Z_i)$
- $Y = \{y\} \cup \bigcup_{i=1}^l V(Y_i)$
- One of $|X|$ or $|Y|$ is odd, so we add Z_1 to that set.

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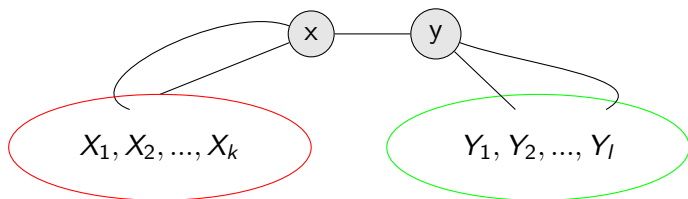


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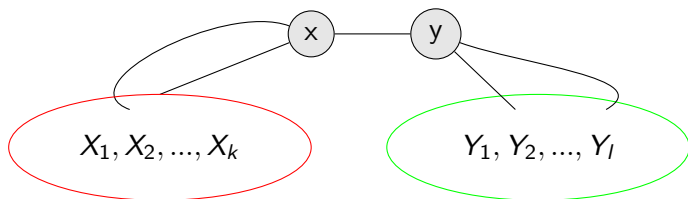
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- $X = \{x, y\} \cup \bigcup_{i=1}^k V(X_i)$
- $Y = \{x, y\} \cup \bigcup_{i=1}^l V(Y_i)$
- Take good partition of $G[X] - W_i$ and $G[Y] - V_i$. If $x \in W_1$, and $y \in V_i$, then $x \in V_1$ and $y \in W_1$, so take union of W_1 and V_1

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A graph G is a basic odd graph if and only if it is a tree with all degrees odd.

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Lemma

A graph G is a basic odd graph if and only if it is a tree with all degrees odd.

Corollary

Let G be a graph, and suppose that every component of G has even order. Then G has a vertex partition such that every vertex class induces a tree with all degrees odd.

Good partition bounds

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$$p(G) = \min \{k : G \text{ has a good partition with } k \text{ sets} \}$$

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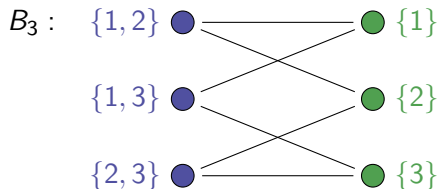
$$p(n) = \max\{p(G) : |G| = n, \text{ every component of } G \text{ has even order}\}$$

Lemma

$$p(n) \geq (1 + o(1))\sqrt{2n}.$$

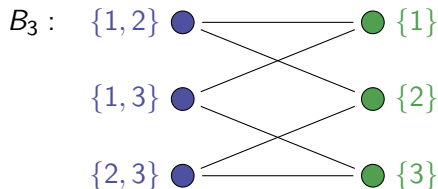
Lemma proof

Let B_n be a biparte graph, with vertex classes $V_0 = [n]^{(1)}$ and $V_1 = [n]^{(2)}$, and edges from $\{i\}$ to $\{j, k\}$ iff $i \in \{j, k\}$.



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$|B_n| = (1 + \frac{1}{n})\frac{n^2}{2}$, so when the order of B_n is even, we can use the first theorem, and find a good partition. Then all singletons have to be in separate components, because if some $\{i\}$, $\{j\}$ were in one group, then $\{i, j\}$ can't be partitioned. Thus $p(n) \geq (1 + o(1))\sqrt{2n}$. after some calculations.

Good partition bounds

Theorem

$$p(n) \leq cn(\log \log n)^{-\frac{1}{2}}.$$

It can be proven that for sufficiently large n , we can find a set W such that $|W| \geq (\log \log n)^{\frac{1}{2}}$, W has all degrees odd, and a good partition. Removing such sets from G leads to upper bound.

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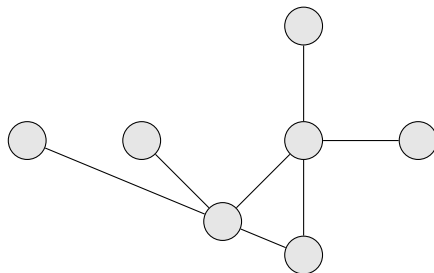
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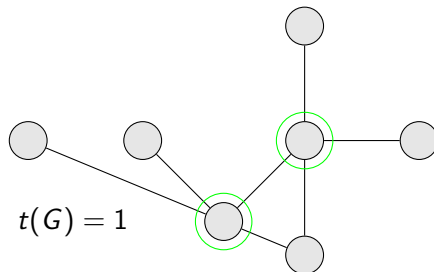
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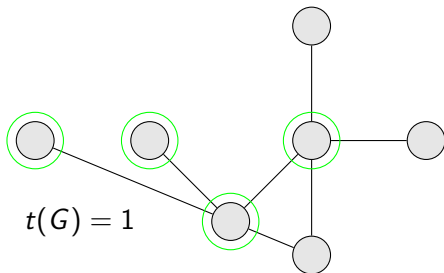
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- If W_i is not an independent set V_{i+1} as the union of the largest set of vertices that induces a subgraph with all degrees odd, of each W_i component
- Use lower bound of $f(n)$ to estimate number of steps in this procedure

Remark: The original proof was made when $f(n) \geq cn$ was not proven yet, so originally the bound was $\leq b \log^2(n)$ with the remark, that if the conjecture is true, then the upper bound is tighter.

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Almost every $G \in \mathcal{G}(n, \frac{1}{2})$, for n even, has a good partition into three sets.

Proof sketch:

Define X_A as an indicator of the event $G[A]$ has all degrees odd

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Using Chybshev inequality bound $n \rightarrow \inf P(X = 0) \rightarrow 0$

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For odd n , almost every $G \in \mathcal{G}(n, \frac{1}{2})$ has a vertex partition $V(G) = V_1 \cup V_2$ such that $G[V_1]$ has all degrees odd and $G[V_2]$ has exactly one vertex with even degree.

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Question

For a random graph $G \in \mathcal{G}(n, \frac{1}{2})$, where n is even, what is $P(p(G) = 2)$

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$f_k(G)$ - the maximum order of an induced subgraph of G with all degrees congruent to 1 modulo k

Definition

$f_k(n) = \min\{f_k(G) : \delta(G) \geq 1 \text{ and } |G| = n\}.$

Lemma

Let $k \geq 2$ be an integer. There exists $c(k) > 0$ such that for every bipartite graph G with $\delta(G) \geq 1$ there is a set $W \subset V(G)$ such that $|W| \geq c(k)|G|$ and $G[W]$ has all degrees congruent to 1 modulo k .

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Theorem

Let $k \geq 2$ be a fixed integer. Let G be a claw-free graph without isolated vertices with order n . Then $f_k(G) \geq (1 + o(1))\sqrt{\frac{n}{12}}$

Previous best result was proven by Y. Caro: $f_k(G) \geq ck(n \log n)^{\frac{1}{3}}$

Claw-graph

